



Revista Facultad de Ingeniería Universidad de Antioquia

ISSN: 0120-6230

ISSN: 2422-2844

Facultad de Ingeniería, Universidad de Antioquia

Díaz, Sergio; Méndez, Diego
Dynamic minimum spanning tree construction and maintenance for Wireless Sensor Networks
Revista Facultad de Ingeniería Universidad de Antioquia, no. 93, 2019, October-December, pp. 57-69
Facultad de Ingeniería, Universidad de Antioquia

DOI: 10.17533/10.17533/udea.redin.20190508

Available in: <http://www.redalyc.org/articulo.oa?id=43062836006>

- How to cite
- Complete issue
- More information about this article
- Journal's webpage in redalyc.org

redalyc.org

Scientific Information System Redalyc
Network of Scientific Journals from Latin America and the Caribbean, Spain and Portugal

Project academic non-profit, developed under the open access initiative

Dynamic minimum spanning tree construction and maintenance for Wireless Sensor Networks

Construcción y mantenimiento dinámico del árbol de expansión mínima para redes inalámbricas de sensores

Sergio Díaz ¹* Diego Mendez ¹

¹Departamento de ingeniería electrónica, Pontificia Universidad Javeriana, Carrera 7 No. 40 - 62, 110231, Bogotá, Colombia.

CITE THIS ARTICLE AS:

S. Díaz, D. Méndez,
"Dynamic minimum
spanning tree construction
and maintenance for
Wireless Sensor
Networks", *Revista Facultad
de Ingeniería Universidad de
Antioquia*, no. 93, pp. 57-69,
Oct-Dec 2019. [Online].
Available:
[https://www.doi.org/
10.17533/10.17533/
udea.redin.20190508](https://www.doi.org/10.17533/10.17533/udea.redin.20190508)

ARTICLE INFO:

Received: October 01, 2018
Accepted: March 18, 2019
Available online: May 07,
2019

KEYWORDS:

Self-organization, topology
control, self-healing

Auto-organización, control
de topología, auto-curación

ABSTRACT: In a Wireless Sensor Network (WSN), finding the optimal route from each node to the sink is not a straightforward task because of the distributed and dynamic characteristics of the network. For instance, the network suffers frequent changes because the channel quality varies over time and the nodes can leave or join the network at any moment. In order to deal with this variability, we propose the Dynamic Gallager-Humblet-Spira (DGHS) algorithm that builds and maintains a minimum spanning tree for distributed and dynamic networks, and we have found that DGHS reduces the number of control messages and the energy consumption, at the cost of a slight increase in the memory size and convergence time. This paper presents a detailed description of the different stages of the DGHS algorithm (neighbor discovery, tree construction and data collection), an in-depth analysis of extensive results that validates our proposal, and an exhaustive description of the GHS limitations.

RESUMEN: En una red inalámbrica de sensores (WSN por su sigla en inglés), encontrar la ruta óptima desde cada nodo al sumidero no es una tarea sencilla debido a las características distribuidas y dinámicas de la red. Por ejemplo, la red sufre cambios frecuentes porque la calidad del canal varía con el tiempo y los nodos pueden abandonar o unirse a la red en cualquier instante. Con el objetivo de controlar esta variabilidad, proponemos el algoritmo dinámico Gallager-Humblet-Spira (DGHS) que construye y mantiene un árbol de expansión mínima para redes dinámicas y distribuidas, y hemos encontrado que DGHS reduce el número de mensajes de control y el consumo de energía, a costa de un ligero aumento en el tamaño de la memoria y el tiempo de convergencia. Este artículo presenta una descripción detallada de las diferentes etapas del algoritmo DGHS (descubrimiento de vecinos, construcción del árbol y recopilación de datos), un análisis profundo de un conjunto extenso de resultados que validan nuestra propuesta, y una descripción exhaustiva de las limitaciones que tiene GHS.

1. Introduction

Wireless Sensor Networks (WSNs) are composed of sensor nodes with constrained resources. The nodes measure a physical variable (e.g., temperature, humidity, or light) in an area of interest, and then, send this information to the sink. WSNs do not rely on pre-existing infrastructure, so, the network topology is built by the nodes once deployed. Initially, the nodes do not have any previous knowledge about their neighborhood because they are scattered randomly, so, they collect information about

their neighbors and the link quality using a procedure called neighbor discovery. Using this information, each node computes routes to reach the sink through multiple hops. In this paper, we propose an algorithm called Dynamic Gallager-Humblet-Spira (DGHS) that builds and maintains a minimum spanning tree which we use to collect data.

In WSNs, finding the optimal route from each node to the sink is not straightforward because of the distributed and dynamic features of the network. On the one hand, the nodes should avoid any centralized control that depletes their energy by sending control packets over long distances. Instead, they should collaborate among them through the exchange of local information. On the other hand, the network is dynamic because of the frequent changes: the channel quality varies over time, and the

* Corresponding author: Sergio Díaz
E-mail: diaz-sergio@javeriana.edu.co
ISSN 0120-6230
e-ISSN 2422-2844

nodes can join or leave the network as a result of mobility or battery exhaustion [1]. Thus, mechanisms for topology construction and maintenance must be distributed and adapt to the dynamic features of the network.

Depending on its internal rules, topology construction mechanisms can lead to different types of structures, such as trees, clusters, or meshes. In the tree topology, each node selects a single parent that receives and forwards the data packets. The parent selection determines how and when the tree is built, so, it is the core of the tree construction mechanisms. The criteria for selecting a neighbor as a parent is different for every application. When the parent selection is over, a tree topology rooted at the sink is available. This tree topology is suitable for WSNs since it enables the many-to-one communication needed to collect data.

Many mechanisms have been proposed to construct a tree topology (See Section 2). However, few mechanisms work in a distributed manner and adapt the tree to the dynamic features of the network. In this paper, we propose a distributed and dynamic algorithm called Dynamic Gallager-Humblert-Spira (DGHS) that builds and maintains a minimum spanning tree. DGHS is based on the Gallager-Humblert-Spira (GHS) [2, 3] algorithm. Note the distinction between DGHS - our mechanism - and GHS - the algorithm proposed by Gallager Humblert Spira. GHS is a static algorithm that finds the minimum spanning tree in a network but does not repair the topology when node failures occur. Additionally, GHS is a theoretical work that has not been evaluated on a wireless network setting. Thus, the contributions of DGHS are twofold: First, DGHS extends GHS by repairing the tree when node failures occur; and secondly, to the best of our knowledge, this is the first paper that evaluates a GHS-based algorithm for a WSN.

DGHS has four phases, namely neighbor discovery, tree construction, data collection, and tree maintenance. In the neighbor discovery, the nodes collect information about their neighbors and the link quality. In DGHS, the link quality is defined as the average number of lost discovery packets. In the tree construction, DGHS finds the minimum spanning tree by executing GHS. In GHS, each node is a different fragment. Subsequently, pairs of fragments merge into a new one. The fragments continue to join until only one remains: the final fragment is the minimum spanning tree. This minimum spanning tree generated by GHS is not rooted at the sink. Thus, in the data collection phase, the sink roots the minimum spanning tree at itself by sending a single control message via the branches of the tree. Then, each node starts sending data packets. In the tree maintenance phase, the nodes repair the tree when communication failures

occur (i.e., the nodes cannot communicate anymore) by merging disconnected fragments. The repair mechanism is initiated by the sink and partially reruns GHS to join disconnected fragments.

We implement DGHS on Contiki which is a widely used open-source operating system for low-power and memory-constrained devices. Subsequently, we evaluate the performance of DGHS on Cooja which is the Contiki network emulator. We compare DGHS with a state-of-the-art tree construction mechanism known as Least Path Interference Beaconing Protocol (LIBP) [4, 5]. The results show that DGHS uses fewer control packets and consumes on average less energy than LIBP, at the cost of a *slight* increase in the memory size and convergence time.

Specifically, the results show that DGHS uses 25.6% fewer control packets than LIBP during the tree construction. Besides, DGHS consumes on average less energy than LIBP: it consumes on average 0.6mW less during the tree construction, and 1.9mW less during the data collection. These results are explained given that DGHS is messageoptimal and sends the data packets through the optimal route (i.e., via the minimum spanning tree). Moreover, DGHS shows a slightly higher memory consumption and slower convergence time because of its complexity.

When compared to our previous work [6], this paper presents: (i) a detailed description of the stages of the DGHS algorithm (neighbor discovery, tree construction and data collection), (ii) an in-depth analysis of extensive results that validates our proposal, and (iii) an exhaustive description of the limitations of GHS for WSNs.

This paper is organized as follows: Section 2 presents recent work about tree construction mechanisms; Section 3 briefly describes GHS and its limitations; Section 4 explains in detail the DGHS algorithm; Section 5 shows the evaluation setting and the performance metrics; Section 6 evaluates the performance of DGHS and LIBP; and, Section 7 presents our conclusions and future work.

2. Related work

The problem of constructing and maintaining a tree topology in a distributed manner is a challenging task. This is because the nodes have limited computational and memory resources and the network is dynamic: the nodes are prone to failures, and the channel quality varies over time [1]. We present relevant work that aims at solving this problem and summarize our findings in Table 1. Note that this table also serves as a qualitative comparison of DGHS with previous work.

Table 1 Previous work regarding tree construction mechanisms

Mechanism	Objective	Evaluation Metrics	Evaluation Setting	Distributed	Dynamic	Limitations
DGHS	Tree construction and maintenance	Number of control packets Energy consumption Convergence time Memory consumption	Emulation	Yes	Yes	A slight increase in memory and convergence time compared to LIBP
GHS [2, 3]	Minimum spanning tree	Communication cost Time complexity	Simulation	Yes	No	See Section 3.2
LIBP[4, 5]	Reduce interfering paths	Power consumption Scalability Throughput Recovery from failure	Emulation	Yes	Yes	It does not assess the network lifetime
TCBDGA[7]	Balance energy consumption	Average tour length Times of reconstructions	Simulation	Yes	Yes	Nodes must know its own location
TST[8]	Minimize tree length	Time complexity Message complexity	Simulation	Yes	No	Nodes must know its own location
HTC[9]	Delay efficient data aggregation	Energy level Response time	Simulation	Yes	Yes	Nodes can have a maximum of two children
Overhear[10]	Secure data collection	Network lifetime Running time	Simulation	No	Yes	It is a centralized approach
RPL[11]	It depends on the objective function	Bandwidth Traffic flow	Testbed	Yes	Yes	It is not appropriate for point-to-point traffic
CTP[12]	Link quality	ETX	Testbed	Yes	Yes	Recovery time
ADCMST[13]	Balance payload	Time complexity Payload	Simulation	No	No	It assumes symmetric links

In [4, 5], the authors present a protocol called Least Path Interference Beaconing (LIBP) that constructs a tree which reduces the interference on the parent nodes. To do that, LIBP limits the number of children that each parent supports. The tree construction is done by sending periodic beacons. Initially, the sink broadcasts a beacon including its identity and weight. Upon reception of this beacon, the nodes reply with an acknowledgment (ACK) informing the sink that it has new children. Then, the sink increments its weight according to the number of new children. By repeating this process, LIBP creates a tree topology that reduces the interference on the parent nodes. Moreover, LIBP repairs the network in the event of parent failure.

[7] presents in-depth an algorithm called TCBDGA (Tree-Cluster-Based Data-Gathering Algorithm). It constructs several tree topologies that gather data using a mobile sink. Initially, the nodes build a single tree topology by selecting a parent according to the residual energy, the distance to the sink, and the local node density. Subsequently, this tree is decomposed into sub-trees considering its depth and traffic load. Finally, the mobile sink visits the sub-trees to collect data and to reconstruct them when their residual energy is low. In [8], the authors present an algorithm called Toward Source Tree (TST) that constructs a tree which forwards information from a source to multiple receivers (i.e., a multicast tree). TST aims at minimizing the tree length and, as a consequence, the delay. To do that, a virtual tree is built by using

only the receivers. Subsequently, pairs of receivers are connected by finding relay nodes between them. Finally, the possible cycles are eliminated from the resultant tree. The drawback of TST is that it assumes that the nodes know their location and it does not work for dynamic networks.

[9] presents an algorithm called Hybrid Tree Construction (HTC) that builds and maintains a tree in which the nodes can only have two children (i.e., a binary tree). By using this tree structure and a time slot scheduling, the nodes aggregate data. The tree construction starts with a neighbor discovery that assesses the energy level and response time of the neighbors. Subsequently, each node selects a parent and a child based on the lowest response time guaranteeing a minimum delay. If two neighbors have the same response time, then the energy level is used as a tiebreaker. Moreover, HTC re-constructs the tree in case of nodes failure.

In [10], the authors construct a shortest-path tree for secure data collection. On the one hand, the tree construction is formulated as (1) an integer linear programming problem; and, (2) a mixed-integer nonlinear programming problem. Both approaches are solved in a centralized manner by the sink. However, the second approach shows a better network lifetime. On the other hand, the secure data collection is implemented by using the overhearing technique. In this technique, a node overhears the incoming and outgoing packets of its

neighbors to determine if they are intentionally dropping or modifying the packets. [11] presents an algorithm called Routing Protocol for Low-Power and Lossy Networks (RPL). It is a multi-hop routing protocol which supports IPv6, being useful for integrating WSNs to the internet. RPL builds a tree topology in which the parent selection is based on a metric referred as *rank*. This metric is proportional to the hop distance from the sink; so, the sink has the lowest rank meanwhile the other nodes increase their rank according to the hop distance. RPL also includes mechanisms for preventing routing loops such as adding a sequence number to the transmitted packets.

The Collection Tree Protocol (CTP) is a data collection mechanism that builds and maintains a tree rooted at the sink [12]. It exchanges beacon messages that transport information about the link quality between neighbor nodes. The node with the highest link quality is selected as next hop, i.e., parent node. CTP has been comprehensively tested on 13 different testbeds, encompassing 7 platforms, and 6 link layers. Besides, it estimates the link quality using information from the physical, link and network layers, expanding the scope.

[13] presents an algorithm called ADCMCST (Algorithm with the minimum number of child nodes) that constructs a tree topology which aims at balancing the nodes payload. To do that, ADCMCST limits the number of children of every parent. Initially, the sink collects the graph of the network and computes an initial tree by using the Breadth First Search algorithm (BFS). Subsequently, the sink limits the number of children by pruning edges on each parent. The algorithm ends if the depth of the tree agrees with a preset value. The drawback of ADCMCST is that it is centralized.

3. Gallager-Humblet-Spira algorithm (GHS)

Since we propose a tree construction and maintenance mechanism based on GHS, we briefly present GHS and its limitations. These limitations are tackled by DGHS.

3.1 Description of the GHS algorithm

GHS [2, 3] finds the minimum spanning tree in a graph with bi-directional edges. It is a distributed algorithm in which the nodes send control messages, wait for a reply, and process the information. In GHS, each node starts as a different fragment. The core idea is to join fragments progressively until there is only one remaining. GHS is implemented by seven control packets, namely connect, initiate, test, accept, reject, report, and change

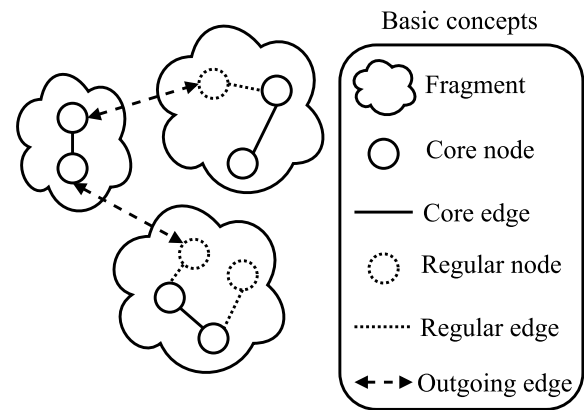


Figure 1 Basic concepts of the GHS algorithm

root. Besides, GHS determines an upper bound on the total number of control messages: for a graph of N nodes and E edges the total number of control messages is at most $5N \log_2 N + 2E$.

Next, we describe the basic concepts of GHS (See Figure 1). A *fragment* is a connected subgraph that belongs to the minimum spanning tree. Fragments can only join via a particular kind of edge, known as the *outgoing edge*. As a general rule, an outgoing edge has a node that belongs to the fragment and another node that does not. Since a fragment could have more than one outgoing edge, the *lowest-weight outgoing edge* of a fragment is an outgoing edge whose weight is the lowest. [2] proves that the lowest-weight outgoing edge of a fragment always belongs to the minimum spanning tree. Additionally, the *core nodes* are the central computing units of the fragment. They find the lowest-weight outgoing edge of a fragment by collecting information about all the outgoing edges. Besides, they initiate the union of two fragments. Finally, the *core edge* is the one that connects two core nodes.

3.2 Limitations of the GHS algorithm for WSNs

We describe the limitations of GHS for a wireless network setting and briefly mention how DGHS deals with those limitations.

- **Packet loss:** GHS does not tolerate packet loss. If for any reason a packet gets lost, GHS will not converge to the minimum spanning tree. The loss of a packet causes that the affected fragment cannot merge and remains disconnected. Thus, it is mandatory to use a reliable packet delivery service. DGHS uses a reliable packet delivery service for the control messages. The reliable packet delivery service is implemented using an acknowledgment packet for every control message. However, this method doubles the number

of control messages. Remember that the total number of control messages is at most $5N \log_2 N + 2E$ for a network without packet loss. So, for a lossy network, such as a WSN, the total number of control messages would increase proportionally to the packet loss rate.

- **Packet errors:** GHS does not tolerate bit errors. If for any reason a packet contains bit errors GHS will not converge to the minimum spanning tree. On the one hand, if the bit error affects a node address, then the packet will be lost. Thus, there will be disconnected fragments, as explained before. On the other hand, if the bit error affects the weight of an edge, then GHS would converge to a non-minimum spanning tree. Thus, to guarantee that DGHS converges to the minimum spanning tree, it is mandatory to use an error detection technique, such as cyclic redundancy check (CRC).
- **Tree rooted at the sink:** The minimum spanning tree generated by GHS is not rooted at the sink. When GHS terminates, the minimum spanning tree is rooted at two core nodes; and, it would be a coincidence if any of these nodes happens to be the sink. DGHS uses an additional control message to root the tree at the sink. It is important to mention that the additional control message does not modify the structure of the minimum spanning tree: it just changes the direction of the edges.
- **Tree maintenance:** GHS finds the minimum spanning tree for a particular network configuration. If the network configuration changes, then the minimum spanning tree is no longer valid. In WSNs, the network configuration is highly dynamic because the channel quality varies over time [1, 14] and the nodes are prone to failures [15, 16]. Thus, as time goes by, the minimum spanning tree loses its validity. As an example, consider the case where the weight of the edges changes in response to an unstable channel. Under that circumstances, the minimum spanning tree is not longer valid because the network configuration has changed. Similarly, a node failure generates a different network configuration. Thus, GHS is not tolerant to network changes since it does not implement any mechanism against them. DGHS copes with the variability of the network by implementing a tree maintenance phase that repairs the tree when node failures occur (i.e., the nodes cannot communicate anymore).
- **Minimum spanning tree:** Finding the minimum spanning tree is a highly energy-consuming process because the nodes require local information about their neighbors and connecting edges, and global information about the different fragments. By

collecting this information, a considerable number of control messages is exchanged. Thus, some tree construction mechanisms [17, 18] do not find the minimum spanning tree that provides the lowest-cost paths. Instead, they find a non-optimal tree whose construction requires fewer control messages but leads to a suboptimal data collection. However, the evaluation results show that this non-optimal tree is more energy inefficient in the long-term compared to the minimum spanning tree.

4. Dynamic Gallager - Humblet - Spira algorithm (DGHS)

The Dynamic Gallager-Humblet-Spira (DGHS) algorithm builds and maintains a minimum spanning tree. To do so, DGHS is divided into four phases, namely neighbor discovery, tree construction, data collection, and tree maintenance (See Figure 2). In the neighbor discovery phase, the nodes collect information about their neighbors and the link quality. DGHS defines the link quality as the average number of lost discovery packets. In the tree construction, DGHS finds the minimum spanning tree by executing GHS. In GHS, each node is a different fragment. Subsequently, pairs of fragments merge into a new one. The fragments continue to join until only one remains: the final fragment is the minimum spanning tree. However, this minimum spanning tree is not rooted at the sink. Hence, in the data collection phase, the sink roots the minimum spanning tree at itself by sending a single control message via the branches of the tree. Then, each node starts sending data packets. In the tree maintenance phase, the nodes repair the tree when communication failures occur by merging disconnected fragments. The repair mechanism is initiated by the sink and partially reruns GHS to join disconnected fragments.

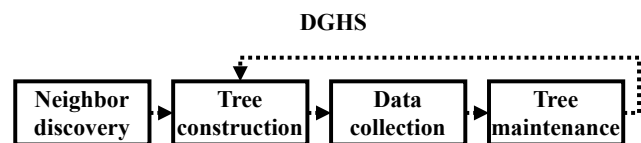


Figure 2 Four phases of the Dynamic Gallager-Humblet-Spira algorithm (DGHS)

4.1 Neighbor discovery

Initially, the nodes do not have any previous knowledge about their neighborhood because they are scattered randomly in the area of interest. So, they collect information about their neighbors and the link quality using a procedure called *neighbor discovery*. During the neighbor discovery, the nodes periodically broadcast

discovery packets with a sequence number. Upon reception of a discovery packet, the nodes compute the link weight which is defined as the average number of discovery packets that have been lost. This computation is done by computing the moving average of sequence number gaps. By exchanging discovery packets, each node maintains a table that includes the link weight for each neighbor.

Specifically, DGHS computes the Exponential Weighted Moving Average (EWMA) of the sequence number gaps. We use EWMA to determine the quality of the links because it applies weights to the samples according to their age. In other words, recent samples strongly influence the resulting average whereas old samples influence the resulting average in a smaller proportion. This property of EWMA is useful in WSNs because it is beneficial to keep an updated neighborhood state which gives more weight to recent samples. We show the definition of EWMA in Equation 1.

$$S_t = \begin{cases} Y_1, & t = 1 \\ \alpha Y_t + (1 - \alpha) S_{t-1}, & t > 1 \end{cases} \quad (1)$$

where, Y_t is the input variable (i.e., in this case the sequence number gaps) and α is a constant between $[0, 1]$ that determines the influence of recent and old samples. If α is close to 1 then the resulting average is almost determined exclusively by recent samples; and, if α is close to 0 then the old samples are given more weight. Once the nodes know their neighborhood, the tree construction phase begins.

4.2 Tree construction

DGHS finds the minimum spanning tree by executing GHS. We select GHS because it is a distributed algorithm, it does not require synchronized nodes, it guarantees an execution without deadlocks, it tolerates an unpredictable delay of messages, and it is message-optimal [3]. Hence, GHS is suitable for WSNs. More importantly, GHS is considered the cornerstone algorithm for finding the minimum spanning tree in a distributed manner. So, several algorithms are based on GHS [19–21]. Note that DGHS inherits the advantages of GHS.

Initially, each node belongs to a different fragment. Subsequently, pairs of fragments merge into a new one. To that end, a fragment identifies its lowest-weight outgoing edge and requests the adjacent fragment to join. After accepting the request, the two fragments join into a new one via the lowest-weight outgoing edge. Pairs of fragments continue to join until only one fragment remains: this final fragment is the minimum spanning tree of the graph. Each fragment has a unique name and a level to guide the joining process. The name guarantees

Algorithm 1 Gallager-Humblet-Spira algorithm - Pseudocode taken from [2] - Part 1

```

procedure wakeup
  begin
    m: Adjacent edge of minimum weight
    SE(m)  $\leftarrow$  Branch
5:   LN  $\leftarrow$  0; SN  $\leftarrow$  Found; Find-count  $\leftarrow$  0
    Send Connect(0) on edge m
  end
procedure Response to Connect(L) on edge j
  begin
10:  if  $L < LN$  then
    begin
      SE(j)  $\leftarrow$  Branch
      Send Initiate(LN, FN, SN) on edge j
      if  $SN = \text{Find}$  then
15:        find-count  $\leftarrow$  find-count + 1
    end
    else if  $SE(j) = \text{Basic}$  then
      place received message on end of queue
    else
20:      Send Initiate(LN+1, w(j), Find) on edge j
    end
procedure Reply to Initiate (L,F,S) on edge j
  begin
    LN  $\leftarrow$  L; FN  $\leftarrow$  F; SN  $\leftarrow$  S; in-branch  $\leftarrow$  j
25:  best-edge  $\leftarrow$  nil; best-wt  $\leftarrow$   $\infty$ 
    for all  $i \neq j$  such that SE(i)=Branch do
      begin
        Send Initiate(L,F,S) on edge i
        if  $S = \text{Find}$  then
30:          find-count  $\leftarrow$  find-count + 1
      end
    if  $S = \text{Find}$  then
      Execute procedure test
    end
35: procedure test
  if there are adjacent edges in the state Basic then
    begin
      test-edge  $\leftarrow$  the minimum-weight adjacent
        edge in state Basic
      Send Test(LN,FN) on test-edge
40:    end
    else
      test-edge  $\leftarrow$  nil
      Execute procedure report
    end
procedure change-root
45:  if SE(best-edge) = Branch then
    Send Change-root on best-edge
  else
    Send Connect(LN) on best-edge
    SE(best-edge)  $\leftarrow$  Branch
50: procedure Response to Change-root
  Execute procedure change-root

```

Algorithm 1 Gallager-Humblet-Spira algorithm - Pseudocode taken from [2] - **Part 2**

```

procedure Response to Test(L,F) on edge j
  begin
    if  $SN = Sleeping$  then
      Execute procedure wakeup
5:   if  $L > LN$  then
      Place received message on end of queue
    else if  $F \neq FN$  then
      Send Accept on edge j
    else
10:   begin
      if  $SE(j) = Basic$  then
         $SE(j) \leftarrow Rejected$ 
      if test-edge  $\neq j$  then
        Send Reject on edge j
15:   else
      Execute procedure test
    end
  end
procedure Response to Accept on edge j
20: begin
  test-edge  $\leftarrow nil$ ;
  if  $w(j) < best-wt$  then
    best-edge  $\leftarrow j$ ; best-wt  $\leftarrow w(j)$ 
  Execute procedure report
25: end
procedure Response to Reject on edge j
  begin
    if  $SE(j) = Basic$  then
       $SE(j) \leftarrow Rejected$ 
30:   Execute procedure test
  end
procedure Report
  if find-count = 0 and test-edge = nil then
    begin
35:    $SN \leftarrow Found$ 
      Send Report(best-wt) on in-branch
    end
  procedure Response to Report(w) on edge j
    if  $j \neq in\_branch$  then
40:   begin
      find-count  $\leftarrow find\_count - 1$ 
      if  $w < best-wt$  then
        best-wt  $\leftarrow w$ ; best-edge  $\leftarrow j$ 
      Execute procedure report
45:   end
    else if  $SN = Find$  then
      Place received message on end of queue
    else if  $w > best-wt$  then
      Execute procedure change-root
50:   else if  $w = best-wt = \infty$  then
      Halt

```

that different fragments merge, and the level determines how and when the joining occurs.

In the remaining of this Subsection, we present in detail the construction of the minimum spanning tree. To that end, we describe the situation where two fragments exchange control messages to join. In other words, we explain a complete round of the tree construction. Note that we included the pseudocode of the minimum spanning tree construction in Algorithm 1 and the nomenclature in Table 2.

Table 2 Nomenclature of GHS

Symbol	Definition
SN	Node State
Sleeping	
Find	Possible values of SN
Found	
SE(j)	State of edge j
Basic	
Branch	Possible values of SE(j)
Rejected	
FN	Fragment identity in the node
LN	Fragment level in the node
F	Fragment identity in the message
L	Fragment level in the message
w(j)	Weight of adjacent edge j
best-edge	Best edge found
best-wt	Weight of the best adjacent edge
test-edge	Edge to be tested
in-branch	Edge that received the initiate msg
find-count	Counter of <i>Find</i> states

Before explaining the details of the GHS algorithm, it is necessary to define the states of the nodes and the edges. A node has two states: find or found. In the find-state, the node is examining its edges to find the one which is outgoing and has the lowest weight. Besides, the node is waiting that its children report their lowest-weight outgoing edge. In the found-state, the node has sent to its parent the information about its lowest-weight outgoing edge. Additionally, the node waits for a request to join its fragment with the adjacent one via its lowest-weight outgoing edge; or, it waits for a message informing that the union was completed by another node within its fragment. An edge has three states: basic, branch, and rejected. A basic edge does not specify whether it belongs to the minimum spanning tree or not; a branch edge belongs to the minimum spanning tree; and, a rejected edge does not belong.

Next, we present the details of the GHS algorithm. In GHS, the nodes send messages, wait for a reply, and process the information. In this way, we describe GHS

by introducing its control messages, namely, connect, initiate, test, accept, reject, report, and change root. In the rest of this Section, we employ two fragments to show the message exchange. Suppose a fragment named fn with a node p and level l ; and, another fragment named fn' with a node q and level l' .

At the very beginning of the algorithm, each node is a fragment with name and level equal to zero. Then, the fragments actively seek to join among them. To that end, each fragment finds its lowest-weight outgoing edge and requests the adjacent fragment [i.e., the fragment on the other side of the lowest-weight outgoing edge] to join. At this moment, finding the lowest-weight outgoing edge of the fragment is trivial because it corresponds to the lowest-weight edge of the single node. This is because fragments are composed of a single node, so, all the edges are outgoing, and there are no other nodes which may have a lower-weight outgoing edge. Moreover, a fragment requests another to join by sending a connect message. Suppose fragment fn requests fragment fn' to join. To that end, node p belonging to the fragment fn sends over the lowest-weight outgoing edge the connect message $\langle connect, l \rangle$. Upon reception of a connect message, node q belonging to the fragment fn' responds depending on the level l' and the status of the edge pq through which the message was received:

- If both fragments have the same level and pq is not a branch edge, then fragment fn' postpones processing the connect message. This connect message is analyzed again when the level l' becomes greater than l , or pq becomes a branch edge.
- If both fragments have the same level and pq is a branch edge, then q sends an initiate message to p . Since both fragments are about to join at an equal level, the new fragment gets a new name and level. The name is equal to the weight of the edge pq , and the level is increased by one. Besides, pq becomes the core edge of the new fragment; p and q become core nodes. Moreover, all the nodes in the new fragment go to the find-state. So, the initiate message is $\langle initiate, weight\ of\ pq, l + 1, find \rangle$.
- If the level l' is greater than l , then q sets the state of the edge qp to branch, and sends an initiate message. Since both fragments are about to join at a different level, the higher-level fragment imposes its name and level over the lower-level fragment. In other words, the lower-level fragment copies the name and level of the higher-level fragment. Thus, the new fragment is called fn' and has level l' . So, the initiate message is $\langle initiate, fn', l', \frac{find}{found} \rangle$. The last parameter is the state of the node q , which can be either find or found, and is represented as $\frac{find}{found}$. This state is assumed by the nodes in the lower-level fragment.

So, if a fragment requests another to join by sending a connect message, it gets an initiate message in response. This initiate message determines the name, level and state of the nodes in the new fragment. Upon reception of the initiate message $\langle initiate, fn, l, \frac{find}{found} \rangle$, a node updates the name and level of its fragment using fn and l , respectively. Besides, the node assumes the state indicated by the last parameter of the message, and selects the sender of the message as parent. Moreover, the node retransmits the initiate message via its branch edges. By retransmitting this message, all the nodes update the name and level of the fragment; and, they also update the route toward the core nodes.

When a node receives an initiate message and the last parameter is find-state, it changes its state to *find*. Immediately after the state change, the node starts searching for its lowest-weight outgoing edge. To do that, the node arranges the edges in a weight-ascending list. Then, it sends a test message to the lowest weight edge. In response, the node could receive an accept message, which means that the edge is outgoing. So, the node has found its lowest-weight outgoing edge, and the search is over. On the other hand, the node could receive a reject message, which means that the edge is not outgoing. In this case, the node marks its lowest-weight edge as rejected, and sends a test message to the next edge in the list. If there are no more edges in the list, the node assigns a weight of ∞ to its lowest-weight outgoing edge. For example, suppose that node p is trying to determine whether node q is outgoing or not. To that end, p sends the test message $\langle test, fn, l \rangle$ to q . Upon reception of the test message, q reacts as follows:

- If the level l is greater than l' , then node q postpones processing the test message. This test message is analyzed again when the level l' has reached or surpassed l .
- If the level l is lower or equal than l' , then q determines whether it belongs to the same fragment than p . To do that, q compares both fragment names. If fn and fn' are equal, then q and p belong to the same fragment. Thus, they cannot join. So, q sends a reject message to p , and marks the edge qp as rejected. On the other hand, if fn and fn' are different, then q and p belong to different fragments. Thus, they can join. So, q sends an accept message to p , and marks the edge qp as lowest-weight outgoing basic edge. In this moment, q has found its lowest-weight outgoing edge: qp .

Once the node has found its lowest-weight outgoing basic edge, it waits for a report message from each child. The report message $\langle report, lwoe \rangle$ contains the lowest-weight outgoing edge (lwoe) that each child knows. By collecting report messages, the node obtains

information about outgoing edges in the fragment and their weights. The node arranges these outgoing edges in a weight-ascending list including its own lowest-weight outgoing basic edge. Subsequently, the node determines that the first element of the list is the lowest-weight outgoing edge. Thus, the node selects the lowest-weight outgoing edge by considering the outgoing edges reported by its children and its own lowest-weight outgoing basic edge. Subsequently, the node informs its parent about its lowest-weight outgoing edge by sending a report message. Once the report message is sent, the node changes its state to *found*.

The report messages propagate in the fragment from the leaf nodes to the core nodes. When a core node has received a report message from each child and has found its lowest-weight outgoing basic edge, it follows the procedure described above to determine its lowest-weight outgoing edge. However, this lowest-weight outgoing edge is fundamental because it corresponds to the lowest-weight outgoing edge of the whole fragment. If the weight of this outgoing edge is ∞ , it means that GHS has found the minimum spanning tree. If the weight is lower than ∞ , the core node sends a changeroot message to node p , which originally reported the outgoing edge pq . The core node knows the route towards p because each node remembers the child that reported the best outgoing edge. Upon reception of a changeroot message, node p changes the state of its lowest-weight outgoing edge to branch, and sends a connect message $\langle connect, l \rangle$ to q .

We are back at the situation where two fragments exchange a connect message to join. Thus, we explained a complete round of the GHS. The tree construction proceeds by joining pairs of fragments until only one remain. It is important to mention that the minimum spanning tree generated by GHS is not necessarily rooted at the sink. Hence, in the next phase, DGHS roots the tree at the sink.

4.3 Data collection

GHS generates a minimum spanning tree that is not rooted at the sink. When GHS terminates, the minimum spanning tree is rooted at two core nodes; and, it would be a coincidence if any of these nodes happens to be the sink. Once GHS terminates, the sink must become the root of the tree. To do that, the sink transmits an initiate message over its branch edges with the parameters set to null. Upon reception of an initiate message, the nodes select the sender as a parent and retransmit the initiate message over their branch edges. By doing so, all the nodes select a new parent that leads to the sink. Thus, the sink becomes the root of the minimum spanning tree. Since the initiate messages are transmitted via the

branch edges, the parent selection does not change the structure of the minimum spanning tree: it just changes the direction of the edges. Once the sink becomes the root of the minimum spanning tree, the data collection starts.

Each node sends a data packet to its parent every X seconds, where X is a random number between $[60, 120]$ seconds. The nodes send data packets using aperiodic intervals of time to prevent nodes from synchronizing and transmitting at the same time, which causes traffic congestion. By sending information via the selected parents, the nodes route the information via the optimal path (i.e., the minimum spanning tree) improving their energy efficiency.

The nodes collect temperature and humidity information using the SHT11 sensor. In this way, the data packet includes 2 bytes to store the temperature, 2 bytes for the humidity, 2 bytes for the source address and 2 bytes for the destination address. Hence, the total packet payload is 8 bytes. We included the data collection phase in this study to measure the energy consumption of the nodes when they sense and transmit temperature and humidity information.

4.4 Tree maintenance

In WSNs, the nodes are prone to failures because of the harsh environments where they are deployed. Node failures generate network partitions leading to uncovered areas. GHS does not implement any mechanism to deal with node failures. In DGHS, we propose a tree maintenance mechanism that repairs the topology when node failures occur.

When one or more nodes suffer a communication failure (i.e., they are unable to transmit and receive any packets) the fragment is divided into two or more fragments: a fragment containing the sink that we call the *main fragment*, and one or more fragments that are disconnected from the main fragment that we call the *sub-fragments*. Our objective is to join again the sub-fragments with the main fragment by partially rerunning GHS. To do so, the sink periodically broadcasts the initiate message $\langle initiate, random\ name, l+1, find \rangle$. Upon reception of this initiate message, the nodes replace their fragment's name with the *random name*, increase their level by one, and assume the state *find*. By assuming the *find* state, the nodes in the main fragment start looking for sub-fragments to join. The *random name* allows that the main fragment and the sub-fragments do not have the same name, so, they can join. Besides, the higher level of the main fragment means that it imposes its name and level when a join occurs. Under this circumstances, the

main fragment merges with the sub-fragments repairing the network. Thus, we have partially rerun GHS to recover the network from a node failure.

DGHS must meet two conditions to rerun GHS successfully. The first condition is that the nodes in the main fragment must wait until the initiate message has arrived at every node in this fragment before start looking for a sub-fragment. This ensures that all the nodes in the main fragment know the *random name* avoiding that they merge among them and generate routing loops. The second condition is that the *random name* must be different from the sub-fragments' name. This ensures that the main fragment and the sub-fragments can recognize among them avoiding that they never join.

5. Performance evaluation setting

We implement DGHS on Contiki which is a widely used open-source operating system for WSNs. Subsequently, we evaluate the performance of DGHS on Cooja which is the Contiki network emulator. We compare DGHS with a state-of-the-art tree construction mechanism known as LIBP (See Section 2). We use LIBP as a point of reference for the comparison because it outperforms well-known tree construction mechanisms, such as RPL [11] and CTP [22], regarding power consumption, scalability, throughput, and recovery from failure [4].

Table 3 Configuration parameters for the emulations

Parameter	Specification
Network emulator	Cooja
Emulated hardware	TelosB node
Operating system	Contiki version 3.x-3029-g2226912
Communication stack	Rime
Radio model	Unit disk graph
Network area	$300 \times 300 \text{ m}^2$
Number of nodes	40
Transmission range	50 meters
Interference range	100 meters
Packet payload	8 bytes
Data packet frequency	[60-120] seconds
α	0.25
Total emulation time	240 seconds

We emulate 40 TelosB nodes in a $300 \times 300 \text{ m}^2$ area. The nodes are randomly scattered. Regarding the wireless channel, we use the unit disk graph model which assumes that the nodes communicate and interfere in fixed-radius circles. Besides, the transmission range is 50 meters, and the interference range is 100 meters. During the data collection, each node sends a data packet every X seconds, where X is a random number between [60,120] seconds;

the packet payload is 8 bytes. We run 20 emulations with different seeds and average the results. Table 3 shows the configuration parameters for the emulations.

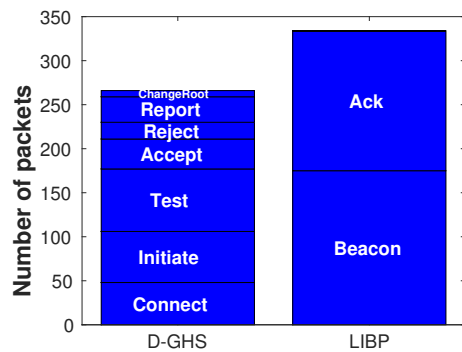
We analyze four metrics to evaluate the performance of DGHS:

- **Number of control packets:** We count the total number of control packets needed to build the tree. Remember that in DGHS the nodes use seven types of control packets to construct the tree (connect, initiate, test, accept, reject, report, and change root); in LIBP, the nodes use two types of control packets (beacon and ACK). Each node counts the number of control packets transmitted, and then we add these values.
- **Energy consumption:** We assess the energy consumption during the tree construction and the data collection. By doing so, we can distinguish between (1) the energy consumption caused by control packets that build the tree, and (2) the energy consumption caused by data packets that traverse the tree. We measure the energy consumption of each node and compute the average.
- **Convergence time:** The convergence time is the number of seconds that the algorithm requires to construct the tree from scratch. So, the convergence time adds the neighbor-discovery time and the tree-construction time. We consider that the tree construction is finished when every node can reach the sink.
- **Memory consumption:** The nodes have constrained memory resources. We measure the percentage of flash and RAM that DGHS consumes on a TelosB node, which has 48KB of flash and 10KB of RAM.

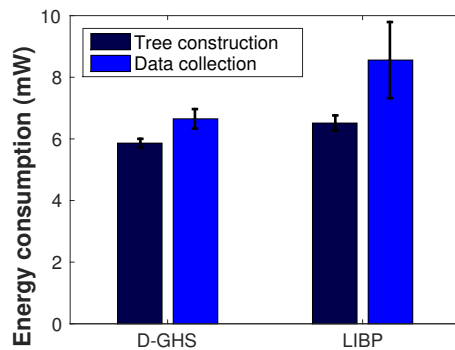
6. Performance evaluation results

We analyze the number of control packets, energy consumption, convergence time, and memory consumption for DGHS and LIBP.

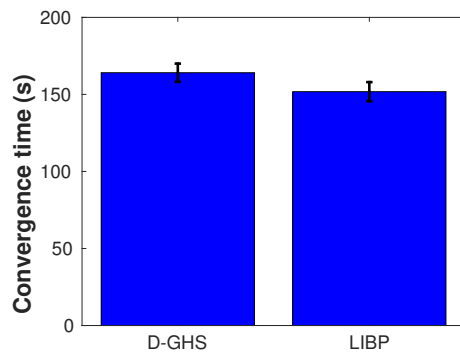
Figure 3a shows the total number of control packets needed to build a tree using DGHS and LIBP. We see that LIBP uses 25.6% more control packets to construct the tree: LIBP employs 334 control packets, and DGHS uses 266. We expected that DGHS uses a lower number of control packets because it is based on GHS which is message-optimal. On the other hand, LIBP constructs the tree by sending periodic beacons and ACKs. This periodic approach does not limit the number of control packets resulting in unnecessary transmissions. Other tree construction mechanisms, such as RPL [11] and CTP [22], implement an adaptive beaconing called the



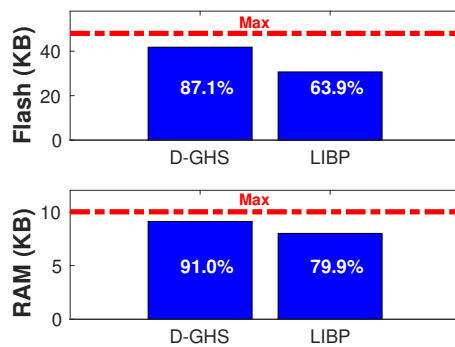
(a) The total number of control packets needed to build a tree



(b) The average energy consumption per node



(c) The convergence time for the tree construction



(d) Percentage of flash and RAM consumed by D-GHS and LIBP

Trickle algorithm [23]. The Trickle algorithm determines the sending rates of beacons in such a way that it sends control packets more often when there are network changes, and it reduces the control traffic rates when the network stabilizes. To do that, the Trickle algorithm uses timers that control the sending rates of beacons. The Trickle algorithm determines the value of the timer by choosing a random value between $[I/2, I]$ where I is initially defined as I_{min} . When the timer expires, RPL sends a beacon and duplicates the value of I up to I_{max} . Additionally, the timer is reset to I_{min} when RPL detects an inconsistency in the network. By setting I_{min} to a few milliseconds and I_{max} to 1 hour, the Trickle algorithm can reduce the control traffic up to 75% compared to sending a beacon every 30 seconds [24].

Figure 3b shows the average energy consumption per node during the tree construction and the data collection. During the tree construction, DGHS consumes 9.2% less energy than LIBP: each node running DGHS consumes 5.9mW on average, and each node running LIBP consumes 6.5mW on average. DGHS consumes less energy than LIBP during the tree construction because it uses fewer control packets. So, the fact that DGHS is based-on a message-optimal algorithm leads to energy savings. During the data collection, DGHS consumes 22.1% less energy than LIBP: each node running DGHS consumes 6.7mW on average, and each node running LIBP consumes 8.6mW on average. DGHS consumes less energy than LIBP because the nodes send the data packets via optimal paths. Remember that DGHS routes the packets via the minimum spanning tree which includes the optimal paths. On the other hand, LIBP routes the data packets via sub-optimal paths causing higher energy consumption. Hence, DGHS is more energy-efficient than LIBP during the tree construction and the data collection. Note that the major energy savings occur in the data collection where the network spends most of the time. So, the long-term energy saving is significant since DGHS could increase the network lifetime by around 22.1%.

Figure 3c shows the convergence time for the tree construction in DGHS and LIBP. We see that DGHS converges 8.1% slower than LIBP: it converges after 164.1 seconds, and LIBP does after 151.8 seconds. The higher convergence time of DGHS is due to the delay of sending information from the outgoing edges to the core nodes using a multi-hop approach. Remember that the core nodes collect information from the whole fragment and determine which is the lowest-weight outgoing edge. So, the higher convergence time of DGHS is because sending information to the core nodes via multiple hops is time-consuming. On the contrary, in LIBP, the nodes rely exclusively on information from their one-hop neighbors. Hence, they avoid sending information over multiple hops

Figure 3 Performance evaluation results

resulting in lower convergence time. In this way, DGHS relies on information provided by neighbors that are N -hops away (note that the value of N is determined by the size of the fragment). On the contrary, LIBP relies on information provided by 1-hop neighbors. So, the delay of packets traversing N -hops increases the convergence time of DGHS. Note that this convergence time is typical of mechanism that gather complete information about the neighborhood by sending packets through N -hops. So, increasing the convergence time is the price paid for having complete information of the neighborhood that results in the optimal paths found by DGHS.

Figure 3d shows the percentage of flash and RAM consumed by DGHS and LIBP on the TelosB nodes. We see that DGHS slightly consumes more memory than LIBP: it consumes 23.2% more flash and 11.1% more RAM. The higher memory consumption of DGHS is because it requires seven types of control messages and LIBP only requires two. The additional control messages of DGHS employ additional structures, lists, and validations causing a higher memory consumption. This result was expected since DGHS can find the minimum spanning tree, so, it is more complex than a protocol that finds a sub-optimal tree such as LIBP. It is important to mention that the flash and RAM usage of DGHS is near to the maximum. DGHS consumes 41.8KB out of 48KB of flash (i.e., 87.1%), and 9.1KB out of 10KB of RAM (i.e., 91.0%). So, DGHS implementation has little space to add new functionalities. A straightforward solution to this problem is to use a node with more memory resources. For example, the RE-Mote (manufacturer Zolertia) has 512KB of flash and 32KB of RAM; the wismote (manufacturer arago-systems) has 256KB of flash and 16KB of RAM. In this way, implementing DGHS in those nodes would leave room for plenty of new functionalities such as adding IPv6.

7. Conclusions and future work

We propose the DGHS algorithm which builds and maintains a minimum spanning tree. To do so, DGHS is divided into four phases, namely neighbor discovery, tree construction, data collection, and tree maintenance. We evaluate the performance of DGHS on Cooja, which is the Contiki network emulator. The results show that DGHS uses 25.6% fewer control packets than LIBP during the tree construction. Besides, DGHS consumes on average less energy than LIBP: it consumes on average 0.6mW less during the tree construction, and 1.9mW less during the data collection. However, DGHS slightly increases the memory size and convergence time because of its higher complexity. In conclusion, DGHS reduces the number of control messages and the energy consumption, at the cost of a slight increase in memory size and convergence time. In future work, we plan to design a fault detection

mechanism that triggers the tree maintenance phase, instead of searching for disconnected sub-fragments periodically. We also plan to evaluate DGHS on a testbed to consider the challenges of a real communication channel.

8. Acknowledgements

The authors would like acknowledge the cooperation of all partners within the *Auto-configurable, self-healing distributed wireless sensor networks supporting high dependability for critical applications* project. The authors would also like to thank all the institutions that supported this work: Pontificia Universidad Javeriana (Project ID: 5685); the German Research Foundation (Deutsche Forschungsgemeinschaft - DFG) (Project ID: KR 3576/21-1), and the Colombian Administrative Department of Science, Technology and Innovation (Colciencias) (Project ID: FP44842-090-2015) through the *Fondo Nacional de Financiamiento para la Ciencia, la Tecnología y la Innovación Francisco José de Caldas*. Also, this work was partially funded by Colciencias through the call 647 of 2014.

References

- [1] A. Castagnetti, A. Pegatoquet, T. N. Le, and M. Auguin, "A joint duty-cycle and transmission power management for energy harvesting wsn," *IEEE Transactions on Industrial Informatics*, vol. 10, no. 2, pp. 928-936, May 2014.
- [2] R. G. Gallager, P. A. Humblet, and P. M. Spira, "A distributed algorithm for minimum-weight spanning trees," *ACM Trans. Program. Lang. Syst.*, vol. 5, no. 1, pp. 66-77, Jan. 1983.
- [3] W. Fokkink, *Distributed Algorithms: An Intuitive Approach*. The MIT Press, 2013.
- [4] L. Ngqakaza and A. Bagula, *Least Path Interference Beaconing Protocol (LIBP): A Frugal Routing Protocol for The Internet-of-Things*. Cham: Springer International Publishing, 2014, pp. 148-161.
- [5] A. Bagula, D. Djenouri, and E. Karbab, "Ubiquitous sensor network management: The least interference beaconing model," in *2013 IEEE 24th Annual International Symposium on Personal, Indoor, and Mobile Radio Communications (PIMRC)*, 2013, pp. 2352-2356.
- [6] S. Díaz, D. Mendez, and M. Scholzel, "Dynamic gallager-humblet-spira algorithm for wireless sensor networks," in *2018 IEEE Colombian Conference on Communications and Computing (COLCOM)*, May 2018, pp. 1-6.
- [7] C. Zhu, S. Wu, G. Han, L. Shu, and H. Wu, "A tree-cluster-based data-gathering algorithm for industrial wsns with a mobile sink," *IEEE Access*, vol. 3, pp. 381-396, 2015.
- [8] H. Gong, L. Fu, X. Fu, L. Zhao, K. Wang, and X. Wang, "Distributed multicast tree construction in wireless sensor networks," *IEEE Transactions on Information Theory*, vol. 63, no. 1, pp. 280-296, Jan. 2017.
- [9] S. Gopikrishnan and P. Priakanth, "Hybrid tree construction for sustainable delay aware data aggregation in wireless sensor networks," *Wirel. Pers. Commun.*, vol. 90, no. 2, pp. 923-945, Sep. 2016.
- [10] W. Y. Alghamdi, H. Wu, W. Zheng, and S. S. Kanhere, "Constructing a shortest path overhearing tree with maximum lifetime in wsns," in *2016 49th Hawaii International Conference on System Sciences (HICSS)*, Jan. 2016, pp. 5858-5867.

- [11] R. Alexander and *et al.*, "RPL: IPv6 Routing Protocol for Low-Power and Lossy Networks," RFC 6550, Mar. 2012. [Online]. Available: <https://rfc-editor.org/rfc/rfc6550.txt>
- [12] T. H. Sardar, A. Khatun, and S. Khan, "Design of energy aware collection tree protocol in wireless sensor network," in *2017 IEEE International Conference on Circuits and Systems (ICCS)*, Dec. 2017, pp. 12–17.
- [13] S. Xie and Y. Wang, "Construction of tree network with limited delivery latency in homogeneous wireless sensor networks," *Wireless Personal Communications*, vol. 78, no. 1, pp. 231–246, Sep. 2014.
- [14] K. Srinivasan, M. A. Kazandjieva, S. Agarwal, and P. Levis, "The beta-factor: Measuring wireless link burstiness," in *Proceedings of the 6th ACM Conference on Embedded Network Sensor Systems*, ser. SenSys '08, New York, USA, 2008, pp. 29–42.
- [15] S. Imran and Y. Ko, "A continuous object boundary detection and tracking scheme for failure-prone sensor networks," *Sensors*, vol. 17, no. 2, 2017.
- [16] S. Chouikhi, I. E. Korbi, Y. Ghamri-Doudane, and L. A. Saidane, "Recovery from simultaneous failures in a large scale wireless sensor network," *Ad. Hoc. Networks*, vol. 67, no. Supplement C, pp. 68–76, 2017.
- [17] Z. Han, J. Wu, J. Zhang, L. Liu, and K. Tian, "A general self-organized tree-based energy-balance routing protocol for wireless sensor network," *IEEE Transactions on Nuclear Science*, vol. 61, no. 2, pp. 732–740, Apr. 2014.
- [18] S. Wan, Y. Zhang, and J. Chen, "On the construction of data aggregation tree with maximizing lifetime in large-scale wireless sensor networks," *IEEE Sensors Journal*, vol. 16, no. 20, pp. 7433–7440, Oct. 2016.
- [19] Y. Choi, M. Khan, V. S. A. Kumar, and G. Pandurangan, "Energy-optimal distributed algorithms for minimum spanning trees," *IEEE Journal on Selected Areas in Communications*, vol. 27, no. 7, pp. 1297–1304, Sep. 2009.
- [20] M. Khan, V. S. A. Kumar, M. V. Marathe, G. Pandurangan, and S. S. Ravi, "Bi-criteria approximation algorithms for power-efficient and low-interference topology control in unreliable ad hoc networks," in *IEEE INFOCOM 2009*, Apr. 2009, pp. 370–378.
- [21] M. Khan and G. Pandurangan, "A fast distributed approximation algorithm for minimum spanning trees," in *Proceedings of the 20th International Conference on Distributed Computing*. Springer-Verlag, 2006, pp. 355–369.
- [22] O. Gnawali, R. Fonseca, K. Jamieson, D. Moss, and P. Levis, "Collection tree protocol," in *Proceedings of the 7th ACM Conference on Embedded Networked Sensor Systems*, New York, NY, 2009, pp. 1–14.
- [23] P. Levis, T. H. Clausen, O. Gnawali, J. Hui, and J. Ko, "The Trickle Algorithm," RFC 6206, Mar. 2011. [Online]. Available: <https://rfc-editor.org/rfc/rfc6206.txt>
- [24] J. Vasseur and A. Dunkels, *Interconnecting Smart Objects with IP: The Next Internet*. San Francisco, CA, USA: Morgan Kaufmann Publishers Inc., 2010.