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Reliability Analysis for Aviation Airline Network Based on Complex Network

Dong Bing^{1,2}

ABSTRACT: In order to improve the reliability of aviation airline network, this paper presents an empirical analysis on the airline network structure of an aviation company in China from the perspective of complex network, and the calculation result of the statistical features and degree distribution of the network, proves that the network is a small-world network and a scale-free network. Four indicators, i.e. degree. closeness, vertex betweenness and flow betweenness, are utilized for aviation network centralization so as to distinguish the most appropriate method. The influence of nodes in local network is to be measured through the indicators. The results show that vertex betweenness can achieve the best aviation network centralization effect. Specifically, the centrality degree reaches 95.87%. On this basis, the network reliability is analyzed to discover that when two nodes with maximum degree or maximum betweenness are removed, the network performance is reduced by a half. Eventually, countermeasures are proposed for further improvement according to the results. In other words, complex network method is feasible used to analyze the topological structure and statistical features of aviation network. Based on this, a study is conduced to the network reliability and suggestions are proposed for optimizing the aviation network.

KEYWORDS: Complex network, Aviation network, Scale-free property, Reliability.

INTRODUCTION

Aviation network refers to an airline system constituted by airlines connected in a certain way in a district, serving as a basis for the production and development of the airline company. In the studies on aviation network with network research methods, it is found that the aviation network has relevant statistical features (Guimera and Amaral, 2004; Guimera et al., 2005; Barrat et al., 2004) of "small-world network". However, most studies (Barrat et al., 2005) focus on the analysis on physical statistical features of aviation network structure and the evolution of overall topological structure, while only a few studies are made on the analysis of route network with social network methods (Porta et al., 2006). Since there are significant differences between aviation network nodes, it is particularly necessary to conduct comparative analyses on relevant nodes and studies on centrality of aviation network. The concept of network centrality can be traced back to the idea of applied statistics in the 19th Century (Gaertler and Wagner, 2001). Typically, different centrality indicators are required for centralization for different types of networks, and the multiple centrality study method needs to be applied in combination with parameters. In China, some scholars adopt other theories and methods to study the centrality of aviation network (Dang and Li, 2011). For instance, some scholars use rank-size model to measure the air transport concentration degree so as to assess the position of hub airports; some scholars mainly adopt the dominant flow method supplemented by squared Euclidean distance method and distance-based cluster method to analyze the level and change of major cities in China in domestic passenger aviation network, based on the air passenger statistical data; moreover, some scholars (Porta et al., 2006) employ the Geographic Information System (GIS) method to study the spatial pattern of the domestic aviation network airport system on the basis of air flow data.

With the development of the civil aviation industry in China, the air transport network keeps booming in scale. Nevertheless, it still suffers imperfection, low reliability and low operation efficiency. On the airline company's side, while planning aviation network, network planners basically conduct decision analysis according to experience. They only take the demand for a single airline as the primary indicator for assessing the necessity of launching an airline. Besides, they often select airlines that are basically the same in route of airlines benefiting other companies while neglecting the network reliability and its overall synergistic effect. The safety and reliability of aviation network exert on an important impact on market competitiveness and economic benefits of an airline company. Therefore, the aviation network should be planned in a systematical manner to improve its overall synergistic effect. A complex network method is proposed to analyze the topological structure and statistical features of aviation network by taking China Southern Airlines (CSA) as an example. Based on this, a study is conduced to the network reliability and suggestions are proposed for optimizing the aviation network of CSA.

COMPLEX NETWORK PROCESSING METHODS AND RELATED RESEARCH

As a small world model and a scale-free network model were proposed in the end of the 20th century, the complex network gradually became the research hotspot in different discipline. In order to facilitate the study of complex network effectively, all kinds of research software are introduced, such as Pajek, Ucinet, NetworkX and NetMiner 3.

In this paper, Ucinet is used for airline network. Ucinet is a social network analysis program developed by Steve Borgatti, Martin Everett and Lin Freeman. The program is distributed by analytic technologies. The software Ucinet involves, in the network analysis, programs such as community discovery and region analysis, ego network analysis and the hole structure analysis and so on. It also contains a large number of analysis programs, such as cluster analysis, multidimensional scaling, singular value decomposition, factor analysis and correspondence analysis, the role and status analysis, including structure, and the role and regular equivalence.

In this paper, we take airline passenger flow data as samples, the city for the network nodes, routes between cities as the network edge, aviation passenger flow between cities as

the mapping relationship between node and nodes in the network structure and construction of air traffic network, as shown in Fig. 1:

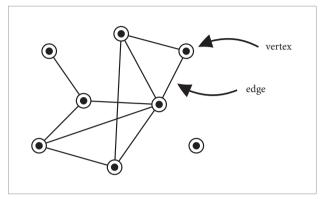


Figure 1. Network structure and construction of air traffic network.

Based on Ucinet software, two kinds of simulation systems were introduced (Hongguang and Liping, 2012), the deliberate targeting system and random interference system are designed, and some simulation experiments are done. The structure diagrams of air passenger flow network are plotted with the software (Dang and Li, 2010) and analyzed from the perspective of structural characteristics, degree of distribution and network centrality. However, traditional research methods fail to thoroughly identify the complexity of aviation network spatial relations between airports in a proper way. Centrality tests for network nodes are important means for judging the importance of nodes in the network, adjusting the aviation layout, and optimizing resource allocation, which, in particular, is greatly significant for safety of aviation network.

EMPIRICAL DATA ACQUISITION AND PROCESSING

In this paper, the domestic and international flight data from CSA's data centre, in 2010, are taken as samples. Usually, cargo flights are arranged at night and the characteristics of transport cargo flights are different from passenger flights. In this article, cargo flights will not be considered. According to the model, let the airport as network node, the direct airline as network edge, and the number of navigable flights among airports as weight of edge, which constituted a weighted aviation network of CSA. Furthermore, adjacency matrix $(K_{ij})_{ij}$ (*n* refers to the number of nodes) is used in order to represent

the aviation network, K_{ij} , in the matrix, refers to the number of flights from Airport i to Airport j. Due to data limitation, the network in this paper is an undirected one. That is, out-degree and in-degree are not involved.

STATISTICAL FEATURE AND THE DEGREE OF DISTRIBUTION

According to the statistics, there are 187 nodes and 1,245 edges. It is thus evident that the aviation network is concentrated and its structure is relatively complex. To show network hierarchical relationship, a backbone aviation network of CSA is built. The node weight threshold is taken as 500 in the weighted network. According to the statistics, there are 99 nodes and 369 edges.

In the complex network theory (Newman, 2003), the indicators reflecting the statistical features of network structure are mainly degree of nodes, average degree, average path length, density, clustering coefficient, etc.

The relation matrix (a_{ij}) is built for the aviation network, in which a_{ij} indicates the relation of flight numbers between city i and city j: if there are any flights between city i and city j, $a_{ij} = 1$; otherwise, $a_{ij} = 0$. Upon calculation, the statistical feature indicators of 2010 aviation network of CSA are listed in Table 1.

Table 1. Statistical features of 2010, CSA aviation network.

	Project	Statistical Value					
1	Number of Nodes	187					
2	Degree of Node	1250					
3	Average Degree	6.684					
4	Average Path Length	2.558					
5	Network Density	0.0358					
6	Clustering Coefficient	0.592					

The degree of a node in an aviation network refers to the number of airports having direct flights with this airport (node); greater degree of a node means greater importance to some extent. Table 2 shows CSA ten airports with top degrees in 2010, among which Guangzhou ranks the first of airport degree as an airline hub of CSA. The average degree of a network refers to the average value of all degrees of nodes in

the network. According to Table 1, the average degree is 6.684, indicating that, in average, each city node is connected to other 6.684 city nodes.

Table 2. Top 10 degree values of CSA airports, in 2010.

	Airport	degree					
1	Guangzhou	106					
2	Shenzhen	48					
3	Beijing Capital	48					
4	Urumqi	46					
5	Changsha	42					
6	Dalian	39					
7	Zhengzhou	38					
8	Shenyang	37					
9	Pudong	34					
10	Wuhan	32					

Generally, the distance between two nodes is defined as the number of edges of the shortest path between two nodes, and the average path length of a network is the average value of the distances between all node pairs. In an aviation network, the distance describes the path from one airport to another one using the minimum transit times; the shorter the distance is, the less transit times are required. Meanwhile, the average path length stands for the depth of the air transport, which is a property of the transport shortcut in the integral network; the shorter the average path length is, the less transit times are required between any two airports, bringing more convenience for the passengers. In 2010, the average path length of CSA aviation network is 2.558, which means only 1.558 transit times are required for transporting from one airport to another, favorably meeting the air transport demand.

An undirected network density is defined as the ratio of actual connection number to the maximum possible connection numbers in the figure. In an aviation network, it describes the ratio of actual number of opened segments to the number of all possible segments. The density indicates the closeness of the air connections among all cities in the network. The value is taken between 0 and 1; if the value is closer to 1, the network structure is more perfect and the connections of the air transport

are closer. In 2010, the density of aviation network is 0.0358, which is relatively small, indicating that the connections are not very strong among CSA airports.

The clustering coefficient of a node in the network stands for the ratio of the actual connection number to the maximum possible connection edges between this node and its adjacent nodes. In an aviation network, the clustering coefficient of a node indicates the average cluster degree of the local network comprised of the airport and its adjacent airports. Higher clustering coefficient means greater cluster degree of the local network, and smaller impact of this node on the adjacent airports; on the contrary, lower value means more dependence of the adjacent airports on this node. Guangzhou Airport has a clustering coefficient of 0.06, which is the smallest of all, indicating that the adjacent airports are highly dependent on Guangzhou Airport, and large numbers of flights will be affected if failure occurs in Guangzhou Airport. The clustering coefficient of the integral network is the average value of the clustering coefficients of all city nodes. As shown in Table 1, CSA aviation network, in 2010, has relatively small average shortest path length and large clustering coefficient; hence this aviation network belongs to a small-world network.

Betweenness is generally defined as the capability of a node for controlling the connection of other node pairs, i.e., the effect of a node acting as a bridge between other node pairs. Higher betweenness of a node indicates stronger effect of the node as a bridge, and more important role in the network. See Table 3 for

Table 3. Betweenness values of top 10 airports of CSA aviation network, in 2010.

	Airport	Betweenness Values					
1	Guangzhou	51.414					
2	Urumqi	25.782					
3	Pudong	9.61					
4	Beijing Capital	8.418					
5	Shenzhen	6.21					
6	Harbin	5.828					
7	Shenyang	5.68					
8	Kunming	4.811					
9	Dalian	4.571					
10	Changsha	4.224					

the betweennesses of top 10 airports of CSA Aviation Network in 2010, based on which Guangzhou Airport, as a CSA hub, has the maximum degree and betweenness, taking up the most important position in the network. Secondly, Urumqi, Beijing Capital and Shenzhen Airports also play prominent roles as bridges. Specifically, Pudong Airport ranks third in the betweenness value; although its degree value is not particularly high, it is still a significant transit in the network. If a fail occurs in important transit airports mentioned above, the connections will be greatly affected between other nodes.

The degree distribution of the nodes in the network can be described in the power-law distribution function;

$$p(k) = k^{-\alpha} \tag{1}$$

The power-law distribution coefficient α shows the degree distribution characteristics of a network(Newman, 2003). The distribution p(k) is the probability of a randomly selected node of degree k.

Power-law distribution is also known as scale-free distribution; a scale-free network is a network whose degree distribution follows a power law. In the log-log plot, a straight line with negative slope can be obtained by conducting a linear fitting for the degree distribution, and the absolute value of slope is the power exponent. If the absolute value is relatively small, this network is scale-free. Correlation coefficient indicates the fitting degree of the curve; higher correlation coefficient means the curve is more favorably fitted, explaining actual problems more adequately. The degree distribution of undirected network is discussed here. By applying a double-segment fitting in the log-log plot, we obtained the power exponent and correlation coefficient variance of the degree distribution (Table 4), and degree distribution (Fig.2) of CSA aviation network in, 2010.

Based on Table 3 and Fig. 2, at the significance level of α =0.0001, Segment 1 and Segment 2 both have excellent fitting, with correlation coefficient exceeding 0.96. Since Segment 1 and Segment 2 both have relatively small power exponents, the degree distribution of CSA aviation network is subject to double-segment power-law distribution, and thus this network is scale-free. Hence, the nodes in the aviation network are heterogeneous, certain nodes (hubs) have large numbers of connections, playing the dominant role in the network, while other large number of nodes only have small numbers of connections and are located on the edge of the network, according to the Matthew Effect.

Table 4. Feature	values o	of degree	distribution	of domestic
flight network.				

	ı	Statistical Value	
1	Number of Nodes(N)		187
2	Degree of Nodes(k)		1250
	Fitting of	Power Exponent($ \alpha $)	0.55094
3	Fitting of Segment 1	Correlation Coefficient(ACC)	0.55094
	Fitting of	Power Exponent($ \alpha $)	1.53157
4	Fitting of Segment 2	Correlation Coefficient(ACC)	0.96051

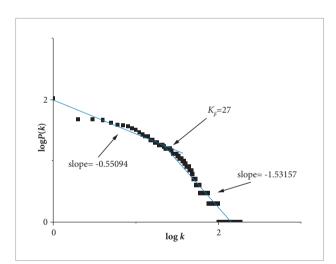


Figure 2. Node degree distribution of CSA aviation network, in 2010.

COMPARISON OF CENTRALITY DEGREE UNDER DIFFERENT INDICATORS

Relevant network centrality indicators serve as the basis for measuring the centrality degree of the network. On the assumption that centrality indicators have been defined in Network C_A with n nodes, the centrality degree of the network will be defined as follows (Costenbader and Valente, 2003):

$$C_A^g = \frac{\sum_{x \in V} (C_A^* - C_A(x))}{(n-1) \max(C_A^* - C_A(x))}$$
 (2)

Where, W refers to the whole network, $C_A^* = \max_{x \in W} C_A(x)$ represents the centrality value of the node with the largest centrality degree.

According to the equation, if the centrality of all nodes is the same, namely the network has no center, then $C_A^* = 0$. In case the centrality degree of only one node is of 1 and that of other nodes is 0, C_A^g will be greater, and the handful of center nodes will be more prominent, which shows that the larger the centrality difference between network nodes, the higher the centrality indicators of the handful of center nodes. Thus, the accuracy of center nodes will be higher, and so will the centrality degree.

Different network centrality (Friedkin, 1991; Newman, 2005) degrees can be figured out by substituting degree, closeness, vertex betweenness and flow betweenness into the following equations respectively.

$$C_{D}^{g} = \frac{\sum_{x \in W} [n - 1 - C_{D}(x)]}{(n - 1)(n - 2)}$$

$$C_{C}^{g} = \frac{\sum_{x \in W} [n - 1 - C_{C}(x)]}{(n - 1)(n - 2)/(2n - 3)}$$

$$C_{B}^{g} = \frac{\sum_{x \in W} [(n - 1)(n - 2) - 2C_{B}(x)]}{(n - 1)^{2}(n - 2)}$$

$$C_{FB}^{g} = \frac{\sum_{x \in W} [n(n - 1) - 4C_{FB}(x)]}{n(n - 1)^{2}}$$
(3)

Where $C_D(x)$, $C_C(x)$, $C_B(x)$, and $C_{FB}(x)$ represents the degree, closeness, vertex betweenness and flow betweenness indicator values of network nodes, respectively. In contrast, the degree indicator is more suitable for measuring the influence of nodes in local network. In global scope, however, the closeness indicator needs to be referenced. The two indicators are only applicable to static network analysis, while the betweenness indicator is more suitable for analysis of dynamic network. When the degree, closeness, vertex betweenness and flow betweenness indicators are used for aviation network centralization based on the formula above, the centrality degree under the different indicators is shown in Table 3.

Obviously, the centrality degree varies with the selected centrality indicators for network centralization. The centrality degree of closeness is relatively low, which indicates that closeness is not suitable for aviation network centralization. Moreover, the top 10 nodes in the aviation network are selected to catch the distribution of indicator values for each node, as shown in Fig. 2. Therefore, degree and closeness fail to obviously distinguish nodes, while the difference of distributions of betweenness is relatively large. Through calculating the centrality degree of each indicator, and comparing the cumulative distribution of centrality data of top 10 nodes, vertex betweenness is the most suitable indicator for aviation network centralization.

NETWORK RELIABILITY ANALYSIS

There are a lot of network attack means in reality (Barrat *et al.*, 2005; Holme *et al.*, 2002; Kai-Quan *et al.*, 2012; Li and Cai, 2012), where the random attack and the hostile attack are relatively representative, and the hostile attack is very destructive. As it was demonstrated before, the aviation network of CSA is a scale-free one, with associate scale-free properties, such as "Stable and Fragile" in attacks, which means it has a very strong flexibility in random attacks or unexpected malfunctions while it is very fragile in hostile attacks (Xiaohuan Wu *et al.*, 2013).

The most important indicators to characterize the network topologic structure are the average path length and the clustering coefficient. The average path length in the aviation network represents air transport depth, the clustering coefficient represents air transport width, and the network efficiency represents overall coordination of the network. With a smaller average path length of the network, a bigger clustering coefficient, and higher network efficiency, it is indicated that the network has a better performance and a stronger fault-tolerant capability. Therefore, this paper will have the reliability analysis on the current network of CSA with the three indicators, and the calculation formulas are as follows:

$$L = \frac{2}{N(N-1)} \sum_{i \ge j} d_{ij} \tag{4}$$

In Eq.(4) N stands for network node number, d_{ij} stands for the shortest distance from Node i to Node j.

$$C_i = 2E_i/(k_i(k_i-1))$$
 (5)

In Eq. (5) k_i stands for the number of edges directly connecting with airport i, and E_i stands for the number of existing connecting edges between airports in number of k_i .

$$E = \frac{1}{N(N-1)} \sum_{i, i \in N, i \neq i} \frac{1}{d_{ii}}$$

$$\tag{6}$$

In Eq.(6) with $0 \le E \le 1$. When E=1, the network is completely connected; and when E=0, all nodes in the network are isolated.

Airport nodes in the aviation network of CSA are sorted according to the degree and the betweenness value in descending order, and then the airport based on a relatively big degree and that, based on a relatively big betweenness, are removed orderly, the average path length, the clustering coefficient, and the network efficiency of the network are calculated respectively, and changes of the average path length, the clustering coefficient and the network efficiency corresponding to the decrease of the airport number are counted and compared, then the contents indicated in Fig. 4, Fig. 5 and Table 5 are respectively obtained.

As indicated by Figs. 3 and 4, the fluctuation of average path length is relatively strong when several nodes are removed, and the fluctuation of average path length based on the remove policy of degree priority is stronger than the one based on the policy of

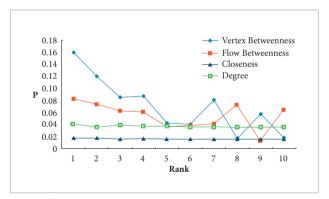


Figure 3. The cumulative distribution of centralized data under four indicators of the top 10 nodes.

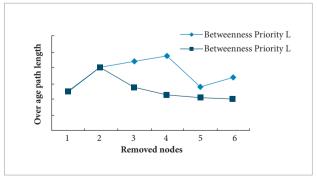


Figure 4. Comparison of average path length changes.

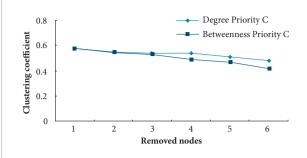


Figure 5. Comparison of clustering coefficient changes.

Table 5. Centrality degree under the different indicators of CSA aviation network.

	Indictor	Centrality Degree					
1	Degree CD	83.12%					
2	Closeness CC	42.50%					
3	Vertex Betweenness CB	95.87%					
4	Flow Betweenness CFB	93.12%					

priority of betweenness. The average path length increases to a largest extent when Guangzhou Airport node with the greatest degree and the highest betweenness is removed, then it decreases afterwards because the network is divided into multiple connected sub-graphs, among which no connection exists. The average path length decreases after increasing, which is a shortcoming for measuring the network performance.

However, overall, it shows that after the node with the greatest degree or the highest betweenness has been removed, the average path length fluctuates widely, which causes the network instability. As it is indicated in Fig. 4, the clustering coefficient decreases after several nodes are removed, and the decrease of the one based on the remove policy of degree priority is slightly larger than the one based on the betweenness priority; therefore, the clustering coefficient decreases and the clustering degree becomes smaller when the node with the greatest degree or highest betweenness has been removed from the network.

As shown in Table 6, when 1~5 nodes with the greatest degree and highest betweenness are removed, the network efficiency suffers a relatively great impact and the backbone network suffers an even greater one, because both of the removed airlines and the flights are of huge numbers in the original network. Especially, when 5 nodes are removed, the efficiency of the original network is decreased by more than 30% and the backbone network decreased by over 70% and the whole network is almost paralyzed. In addition, the decrease based on the remove policy of degree priority is larger than that based on the betweenness priority. In order to understand the network change condition deeply, taking the test of removing two cities one-time for example, we conduct specific analysis on the network performances, based on the two kinds of remove policies before and after the test. Before the test, backbone network of the CSA

Table 6. Comparison of network efficiency changes based on two different remove policies.

	Based on Degree Priority						Based on Betweenness Priority							
Test Procedure	Original Network		Backbone Network		Impacted	Original Network		Backbone Network			Impacted			
	Node	Edge	Efficiency	Node	Edge	Efficiency	Flights	Node	Edge	Efficiency	Node	Edge	Efficiency	Flights
Before Test	187	1245	_	99	369	_	_	187	1245	_	99	369	_	_
Remove 1 Point	186	1035	-11.21	98	267	-34.68	18.6	186	1035	-10.35	98	267	-34.68	18.6
Remove 2 Points	185	939	-18.36	97	228	-51.46	28.4	185	945	-17.54	97	237	-48.73	25.42
Remove 3 Points	184	847	-25.50	96	188	-60.78	37.9	184	882	-23.26	96	217	-56.83	34.26
Remove 4 Points	183	761	-30.64	95	160	-69.47	40.1	183	790	-29.01	95	177	-64.30	39.62
Remove 5 Points	182	685	-33.73	94	136	-76. 31	43.3	182	700	-31.43	94	142	-72.46	40.05

contains 99 nodes and 369 routes; when two airports with the greatest degree value (Guangzhou Airport and Shenzhen Airport) have been removed according to the first policy, 306 routes disappear, the original network efficiency is decreased by 18.36%, the backbone network efficiency is decreased by 51.46%, and the network performance is decreased by half. When two airports with the highest betweenness (Guangzhou Airport and Urumqi Airport) are removed according to the second policy, 300 routes disappear, the original network efficiency is decreased by 17.54%, and the backbone network efficiency is decreased by 48.73%. Its decrease is slightly smaller than that of the remove policy based on the degree priority.

When two airports with the greatest degree are removed, the number of routes of the whole network in south-to-north direction is decreased substantially; when two airports with the highest betweenness are removed, most airports in the western area become isolated nodes, holding up flights, and the overall network performance is decreased by half or so. However, the network performance decreases approximately by about 25% (Dang and Li, 2011) when two nodes are removed in the America aviation network, indicating that the southern aviation network is relatively fragile when facing a selective attack and the overall network reliability is expected to be enhanced. Apart from the selective attack, the current network layout of CSA can cause heavy strike to the aviation transportation when facing random attacks, including natural disasters. Therefore, while enlarging the scale of network development, CSA should also develop its aviation network into a multihub system. Judging from the degree value and betweenness value of the airport node, number of airlines removed from the network, and the network efficiency decrease condition, in addition to taking Guangzhou Airport as the core hub, CSA can further plan and revise its aviation network by defining Beijing as the important hub between Europe, America and the inland of China, defining Urumqi as the regional hub between the middle Asia and the inland of China. On the other hand, the Company can promote connections in east-to-west direction, for example, establishing connections between the west and the east centered in Zhengzhou to fill up the blank in this direction. In this way, the overall performance and reliability of the network can be enhanced, ensuring fluent running of the air transport system.

CONCLUSION

In this paper, complex network theory is applied, an aviation network structure model is built for CSA, and its structure is analyzed. It is discovered that Guangzhou Baiyun Airport is the hub of the aviation network of CSA, the majority of airlines are in the south and north directions, the airlines in Western Region are distributed around Urumqi in a radial way. Moreover, analysis is conducted on the distribution of statistical features and degree distribution of the network, and it is proven that the aviation network of CSA is a small-world network and a scale-free network. Based on this, the reliability of the network is analyzed according to the degree and betweenness of nodes in the network. The results show that the backbone network performance will be reduced to a half, once 2 nodes are removed and it basically breaks down once 5 nodes are removed. The overall reliability of the network is far from high. Therefore, CSA should put more efforts in the overall programming of the aviation network as well as the construction and management of the aviation hub, seek to develop as an multi-hub airport, launch more flights in west and east directions, reasonably allocate air transport resources, and improve the overall performance of the network so as to meet the demand of sustained and healthy development of air transport of CSA.

The Complex network method is used to analyze the topological structure and statistical features of aviation network. Based on this, the suggestions are proposed for optimizing the aviation network from the analysis of the network reliability the application of the methodology (airport planning, fleet sizing, route planning, etc) isn't mentioned in the paper. They will be discussed in other papers.

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